

1 Program Correctness: Template

Given an array b of type **integer** \rightarrow **boolean** and a non-negative length n . The following program computes the numerical representation of a bitvector.

```

i := 0;
x := 0;
while i < n do
  x := 2 · x;
  if  $b[i]$  then
    x := x + 1
  fi;
  i := i + 1
od

```

Let S be above program. Prove the following partial correctness formula:

$$\{n \geq 0\} S \{x = \sum_{i=0}^{n-1} 2^{n-1-i} \cdot (b[i] ? 1 : 0)\}$$

Here is a proof outline in proof system PW:

```

{n ≥ 0}
i := 0;
x := 0;
{inv : i ≤ n ∧  $x = \sum_{j=0}^{i-1} 2^{i-1-j} \cdot (b[j] ? 1 : 0)$ }
while i < n do
  {i < n ∧  $x = \sum_{j=0}^{i-1} 2^{i-1-j} \cdot (b[j] ? 1 : 0)$ }
  {i < n ∧ 2 · x = 2 ·  $\sum_{j=0}^{i-1} 2^{i-1-j} \cdot (b[j] ? 1 : 0)$ }
  {i < n ∧ 2 · x =  $\sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
  x := 2 · x;
  {i < n ∧  $x = \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
  if  $b[i]$  then
    {i < n ∧  $x = \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0) \wedge b[i]$ }
    {i < n ∧  $x + 1 = (b[i] ? 1 : 0) + \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
    x := x + 1
    {i < n ∧  $x = (b[i] ? 1 : 0) + \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
  else
    {i < n ∧  $x = \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0) \wedge \neg b[i]$ }
    {i < n ∧  $x = (b[i] ? 1 : 0) + \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
  skip
  {i < n ∧  $x = (b[i] ? 1 : 0) + \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
  fi;
  {i < n ∧  $x = (b[i] ? 1 : 0) + \sum_{j=0}^{i-1} 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
  {i < n ∧  $x = \sum_{j=0}^i 2^{i-j} \cdot (b[j] ? 1 : 0)$ }
  i := i + 1
  {i ≤ n ∧  $x = \sum_{j=0}^{i-1} 2^{i-1-j} \cdot (b[j] ? 1 : 0)$ }
od
{i = n ∧  $x = \sum_{j=0}^{i-1} 2^{i-1-j} \cdot (b[j] ? 1 : 0)$ }
{x =  $\sum_{i=0}^{n-1} 2^{n-1-i} \cdot (b[i] ? 1 : 0)$ }

```